## On Tree-based Methods for Similarity Learning

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## Outline

#### Introduction

TreeRank for bipartite ranking

Similarity TreeRank

# Biometric verification (1/2)

## A biometric system uses:

► Two **measurements** *X* and *X*′,



- A **similarity** *S* that quantifies the likeness of (X, X'),
- A **threshold** *t* that separates positive and negative pairs.

Aim: S(X, X') > t is a good indicator of Z = +1 with:

 $Z = \begin{cases} +1 & \text{if } (X, X') \text{ from the same person,} \\ -1 & \text{otherwise.} \end{cases}$ 

Two types of errors:

$$\mathsf{TPR}_{S}(t) := \mathbb{P}\{S(X, X') > t \mid Z = +1\}, \\ \mathsf{FPR}_{S}(t) := \mathbb{P}\{S(X, X') > t \mid Z = -1\}.$$

The set  $\{(FPR_{S}(t), TPR_{S}(t)) \mid t \in \mathbb{R}\}$  is known as the **ROC curve**.

# Biometric verification (2/2)

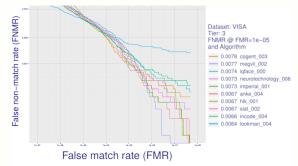


Figure: Extract of the NIST Face Recognition Vendor Test (FRVT) report.

Several criterions measure ROC accuracy:

- Area Under the ROC Curve (AUC),
- Pointwise ROC optimization (pROC) see [Vogel et al., 2018],
- Local AUC (LocAUC): see [Clémençon and Vayatis, 2007].

#### **Remarks:**

- · AUC does not focuses on best instances,
- **pROC** and **LocAUC** can have inadapted optimums.

# Our contribution

The posterior probability  $\eta(x, x') = \mathbb{P}\{Z = +1 | (X, X') = (x, x')\}$ , is the **optimal similarity** in a ranking and ROC sense.

### **Contributions:**

- Procedure for building a similarity that approximates η.
- Guarantees in **sup norm** in the **ROC space**.
- Implementations in specific cases.

## Plan:

- 1. Present **TreeRank for bipartite ranking**, see [Clemencon and Vayatis, 2009].
- 2. Adapt it to similarity learning.
- 3. Draw on **U-statistic theory** to prove new results.

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## **Bipartite ranking**

Standard binary classification framework:

- let a random variable  $(X, Y) \in \mathcal{X} \times \{-1, +1\},\$
- *n* i.i.d. copies  $\mathcal{D}_n = \{(X_i, Y_i)\}_{i=1}^n$  of (X, Y),
- the distribution of (X, Y) is summarized by  $(\mu, \eta)$  where:

$$\forall x \in \mathcal{X}, \ \eta(x) = \mathbb{P}\{Y = +1 | X = x\}, \ \mu(C) = \mathbb{P}\{X \in C\}.$$

• or by 
$$(p, \alpha, \beta)$$
, where  $p = \mathbb{P}\{Y = +1\} = 1 - q$ ,  
 $\alpha(C) = \mathbb{P}\{X \in C | Y = -1\}$  and  $\beta(C) = \mathbb{P}\{X \in C | Y = +1\}$ ,

 $\alpha$  the false positive rate (FPR),  $\beta$  the true positive rate (TPR).

**Objective:** Rank items in  $\mathcal{X}$  by decreasing  $\eta$  using  $\mathcal{D}_n$ .

The ranking is derived from **a scorer**  $s : \mathcal{X} \to \mathbb{R}$  that ranks  $\mathcal{X}$  with  $\mathbb{R}$ .

**Optimal** scorers  $S^*$  are **increasing transforms**  $T \circ \eta$  of  $\eta$ . A scorer  $s^*$  is optimal i.f.f.  $\exists w \in \mathcal{L}_1, w \ge 0$  and V cont. r.v. in (0, 1):

$$\forall x \in \mathcal{X}, \quad s^*(x) = \inf_{z \in \mathcal{X}} s^*(z) + \mathbb{E}\left[w(V) \cdot \mathbb{I}\left\{\eta(x) > V\right\}\right].$$

An approach to bipartite ranking (1/2)

Idea: Using estimated super-level sets of  $\eta$ :  $\{x \in \mathcal{X} | \eta(x) > t\}_{t \in \mathbb{R}}$ , build an accurate scorer.

Optimal scorers write as:  $s^*(x) = C + \mathbb{E}[w(V) \cdot \mathbb{I}\{\eta(x) > V\}]$ , which can be **approximated by a piecewise constant function**,

$$s_N(x) = \sum_{j=1}^N \mathbb{I}\{x \in R_j\},$$

with  $\{R_j\}_{1=1}^N$  increasing  $(R_j \subset R_{j+1})$  family of sets.

The ROC curve of  $s_N$  is the broken line connecting the dots:

$$\{(\alpha(R_j), \beta(R_j))\}_{0 \le j \le N}. \text{ with } R_0 = \emptyset.$$
(1)

Hence, if the  $R_j$ 's are the level sets of  $\eta$ , eq. (1) could be a **piecewise linear** approximation of the ROC curve.

## An approach to bipartite ranking (2/2)

From the **Neymann-Pearson fundamental lemma**, the optimal solution of pROC at level  $\alpha$ , i.e.

$$\max_{\mathcal{C}} \beta(\mathcal{C}) \text{ s.t. } \alpha(\mathcal{C}) \le \alpha, \tag{2}$$

is  $\{x \in \mathcal{X} | \eta(x) > \gamma\}$  where  $\gamma$  is the  $(1 - \alpha)$ -quantile of  $\eta(X) | Y = -1$ . [Clémençon and Vayatis, 2009] solves eq. (2) to build good scorers.

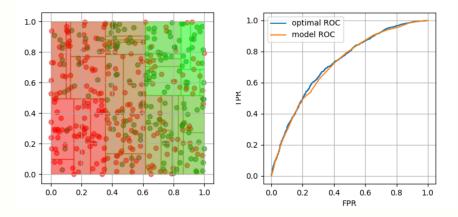
The weighted classif loss:  $L_c(C) = cp \cdot \beta(C) - (1 - c)q \cdot \alpha(C)$ , has for optimal solution  $C = \{x \in \mathcal{X} | \eta(x) > 1 - c\}$ .

TreeRank [Clemencon and Vayatis, 2009] exploits that idea.

TreeRank splits  $\mathcal{X}$  recursively to retrieve the super-level sets of  $\eta$ , with weighted classif. optimized on a family  $\mathcal{C}$  of subsets of  $\mathcal{X}$ .

## Visual example of TreeRank

For C being **coordinate splits**,  $\mathcal{X} = [0, 1]^2$  and  $\eta(x) = (4x_1 + 2x_2)/7$ :



**Remark:** Elements of C are very different from super-level sets of  $\eta$ .

# TreeRank algorithm

**Input.** Maximal depth D > 1 of the tree,  $C, D_n$ . 1. (INIT.) Set  $C_0 = \mathcal{X}$ ,  $\alpha_{d,0} = \beta_{d,0} = 0$  and  $\alpha_{d,2^d} = \beta_{d,2^d} = 1$  for all  $d \ge 0$ . 2. (ITERATIONS.) For d = 0, ..., D - 1 and  $k = 0, ..., 2^{d} - 1$ : 2.1 (OPTIMIZATION STEP.) Set the entropic measure:  $\widehat{\Lambda}_{d\,k+1}(\mathcal{C}) = (\alpha_{d\,k+1} - \alpha_{d,k}) \cdot \widehat{\beta}(\mathcal{C}) - (\beta_{d,k+1} - \beta_{d,k})\widehat{\alpha}(\mathcal{C}).$ Find  $C_{d+1,2k} \in C$  subset of  $C_{d,k}$  that maximizes  $\Lambda_{d,k+1}$ . 2.2 (UPDATE.) Set  $\alpha_{d+1,2k+1}$ ,  $\beta_{d+1,2k+1}$ ,  $\alpha_{d+1,2k+2}$  and  $\beta_{d+1,2k+2}$ . 3. (OUTPUT.) After *D* iterations, get the piecewise constant score function:  $s_D(x) = \sum (2^D - k) \mathbb{I}\{x \in \mathcal{C}_{D,k}\},\$  $\beta_{1,1}$  · · / First split • For pruning: [Clemencon and Vayatis, 2009]. Score has · For bagging: [Clémençon et al., 2013]. TPR 2 values FPR  $\alpha_{1.1}$ 

## Guarantees

## **Assumptions:** Let ≺ denote **absolute continuity**, <sup>1</sup>

- **1.**  $\alpha$  and  $\beta$  are **equivalent** (i.e.  $\alpha \prec \beta$  and  $\beta \prec \alpha$ ),  $\frac{d\beta}{d\alpha}$  is bounded.
- **2.**  $\eta(X) \prec \lambda$  with  $\lambda$  the Lebesgue measure.
- **3.** C is of VC-dimension V,
- **4.** contains all level sets of  $\eta$  and is intersection-stable ( $C \cap C' \in C$ ),

**Theorem:** Under those, given a tree of depth  $D = D_n \sim \log(\sqrt{n})$ ,  $\forall \delta > 0, \exists \lambda > 0$  such that, with probability  $\geq 1 - \delta, \forall n \in \mathbb{N}$ ,

$$\|\operatorname{ROC}_{s_{\mathcal{D}_n}} - \operatorname{ROC}_{s^*}\|_{\infty} \le \exp(-\lambda \sqrt{\log(n)}).$$

 $\lambda$  depends on V,  $\delta$  and universal constants.

 $^{1}\mu \prec \nu \text{ i.f.f. } \exists h : \mathcal{X} \rightarrow \mathbb{R}^{+}, \mu = h \cdot \nu.$ 

## Sketch of proof ([Clemencon and Vayatis, 2009]) Step 1: Bound the $\|\cdot\|_{\infty}$

With  $\beta_{D,k}^*$ ,  $\alpha_{D,k}^*$  params for the *adaptative broken line* est. of ROC<sub>s\*</sub>, when  $t \in [\alpha_{D,k_0}^*, \alpha_{D,k_0+1}] \cap [\alpha(R_{D,k-1}), \alpha(R_{D,k})]$ ,

$$\begin{aligned} \mathsf{ROC}_{\mathsf{s}^*}(t) &\leq \beta^*_{\mathcal{D},k_0} + \mathsf{ROC}'_{\mathsf{s}^*}(0) \times (t - \alpha^*_{\mathcal{D},k_0}), \\ \mathsf{ROC}_{\mathsf{s}_{\mathcal{D}}}(t) &\geq \beta(\mathsf{R}_{\mathcal{D},k}) - \mathsf{ROC}'_{\mathsf{s}^*}(0) \times (t - \alpha(\mathsf{R}_{\mathcal{D},k})). \end{aligned}$$



It implies the following bound on  $\|ROC_{s^*} - ROC_{s_p}\|_{\infty}$ :

$$\max_{1 \le k \le 2^{D}-1} \beta_{D,k}^{*} - \beta(R_{D,k-1}) + \text{ROC}'_{s^{*}}(0) \left[ \alpha_{D,k}^{*} - \alpha(R_{D,k-1}) \right].$$
(3)

**Step 2: Prove by recurrence a bound of eq. (3)** Under assumptions,  $\exists K$  s.t.  $\forall \delta > 0$ , with probability  $1 - \delta$ ,  $\forall d, k$ :

$$|\alpha_{d,k}^* - \alpha(R_{d,k-1})| + |\beta_{d,k}^* - \beta(R_{d,k-1})| \le K^d B(d+1,n,\delta),$$

where  $B(d + 1, n, \delta) = O\left(\frac{V^{1/2d} + \log(1/\delta)^{1/2d}}{n^{1/2d}}\right)$ . Which is proven using **standard VC inequalities**.

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# Similarity ranking

We chose the standard **classification** framework:

- let a random variable  $(X, Y) \in \mathcal{X} \times \{1, \dots, K\}$ ,
- the distribution of (X, Y) is summarized by  $(\mu, (\eta_1, \ldots, \eta_K))$ :

 $\forall x \in X, k \in \{1, \ldots, K\}, \quad \eta_k(x) = \mathbb{P}\{Y = k | X = x\}.$ 

• the optimal similarity is  $\eta(x, x') = \sum_{k=1}^{K} \eta_k(x) \eta_k(x')$ , i.e. the probability to be in the same class.

**Objective:** Rank pairs in  $\mathcal{X} \times \mathcal{X}$  by decreasing  $\eta$  using  $\mathcal{D}_n$ .

The ranking is derived from **a similarity**  $s : \mathcal{X} \times \mathcal{X} \to \mathbb{R}$ .

Given two i.i.d. pairs (X, Y) and (X', Y'), set  $Z = 2\mathbb{I}\{Y = Y'\} - 1$ , one can form n(n-1)/2 obs. of the form ((X, X'), Z) from  $\mathcal{D}_n$ .

Idea: Run TreeRank on non-i.i.d. data of the form ((X, X'), Z).

## TreeRank for similarity ranking

**TreeRank** on data of the form ((X, X'), Z) gives a similarity s.

Similarities satisfies more constraints than scorers.

**Symmetricity:** For *s* to be symmetric, it suffices that C is symmetric.  $\rightarrow$  Use symmetric proposal regions,

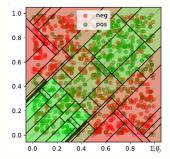
 $\rightarrow$  Learn on data of the form ((X + X', |X - X'|), Z).

Lemma: Let  $s : \mathcal{X}^2 \to \mathbb{R}$ , s is symmetric  $\Leftrightarrow \exists s_0 : \mathcal{X}^2 \to \mathbb{R}$ ,  $s(x, x') = s_0(x + x', |x - x'|)$ .

#### **Identity:**

One expects  $s(x,x) \ge s(x,z)$  for all  $x, z \in \mathcal{X}$ . It is not satisfied by:

$$\eta(x,x') = \sum_{k=1}^{K} \eta_k(x) \eta_k(x').$$



## Extension of the results

Same type guarantees hold for similarity TreeRank.

Estimates of  $\alpha$ ,  $\beta$  are ratios of *U*-statistics (means of pairs), e.g.:

$$\hat{\alpha}(C) = \frac{1}{n_{-}} \sum_{i < j} \mathbb{I}\{Z_{i,j} = -1, (X_i, X_j) \in C\},\$$

#### and standard VC inequalities do not hold.

Using the fact that a simple U-statistic  $U_n(h)$  can be rewritten as:

$$U_n(h) := \frac{2}{n(n-1)} \sum_{i < j} h(X_i, X_j) = \frac{1}{n!} \sum_{\sigma \in \mathfrak{S}_n} \frac{1}{\lfloor n/2 \rfloor} \sum_{i=1}^{\lfloor n/2 \rfloor} h(X_{\sigma(i)}, X_{\lfloor n/2 \rfloor + i}),$$

which is the **first Hoeffding decomposition**, Jensen's inequality implies new VC inequalities, see [Vogel et al., 2018].

Bagging TreeRank models learned on **incomplete U-stats** works, see [Clémençon et al., 2016].

## Conclusion

#### In a nutshell:

TreeRank describes a general approach to the ranking problem.

Performance of **similarity learning** algorithms is evaluated by **ranking criterions** in important applications.

TreeRank can be adapted for **similarity learning**, and theoretical results extended **using results on** *U***-statistics**, see [Clémençon et al., 2016].

TreeRank's competitiveness depends on **the expressivity of** C.

#### **Future work:**

Explore the idea of **using neural networks** to represent C.

# Merci!

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